# On Latin Young Diagrams <br> Daniel Kotlar <br> Tel-Hai College 

Joint work with Ron Aharoni, Eli Berger and He Guo

Young diagram

$$
Y=
$$



## Young diagram


$m$ rows: $\quad r_{1}, r_{2}, \ldots, r_{m}$
row lengths: $a_{1}, a_{2}, \ldots, a_{m}$ respectively
$a_{1}$ columns: $c_{1}, c_{2}, \ldots, c_{a_{1}}$

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|Y|=\sum_{i=1}^{m} a_{i}
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# Question (Chow, Fan, Goemans, Vondrak 2003*): Is it possible to place in each row $r_{i}$ the numbers $1,2, \ldots, a_{i}$ so that the entries in each column are distinct? 

[^0]
## Young diagram


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## Question (Chow, Fan, Goemans, Vondrak 2003*):

 Is it possible to place in each row $r_{i}$ the numbers $1,2, \ldots, a_{i}$ so that the entries in each column are distinct?
## If the answer is Yes we say that $Y$ is Latin

[^1]
## Some easy examples

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a_{1}=a_{2}=\ldots=a_{m}=k
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& m \leq k
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| :--- | :--- | :--- |
| 2 | 1 |  |
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|  |  |  |



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| :--- | :--- | :--- |
| 2 | 1 |  |
| 1 | 2 |  |

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| :--- | :--- | :--- |
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| :--- | :--- | :--- |
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| 1 | 2 |  |
|  |  |  |

No

| 2 | 3 | 1 |
| :--- | :--- | :--- |
| 1 |  |  |
|  |  |  |


| 3 | 2 |
| :--- | :--- |

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|  |  |  |
|  |  |  |

No

| 2 | 3 | 1 |
| :--- | :--- | :--- |
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| 2 | 1 |  |
|  |  |  |
|  |  |  |
|  |  |  |


| 6 | 5 | 4 | 3 | 2 | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 4 | 3 | 2 | 1 |  |  |
| 3 | 2 | 1 |  |  |  |
| 1 |  |  |  |  |  |


|  |  |  |  |
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| :---: | :---: | :---: | :---: | :---: | :---: |
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Yes

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| :--- | :--- | :--- | :--- |
| 1 | 3 | 2 |  |
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No

These diagrams are "too narrow"

Wide diagrams

## Wide diagrams

$$
\begin{aligned}
& A=\left(a_{1}, a_{2}, \cdots\right), a_{1}+a_{2}+\cdots=n, \quad a_{1} \geq a_{2} \geq \cdots \\
& B=\left(b_{1}, b_{2}, \ldots\right), b_{1}+b_{2}+\cdots=n, \quad b_{1} \geq b_{2} \geq \cdots \\
& A \text { dominates } B \text { if } \quad \forall k \quad \sum_{i=1}^{k} a_{i} \geq \sum_{i=1}^{k} b_{i}
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Definition (Chow et al.):
A Young diagram $Y$ is wide if every sub-diagram $Z$ formed by a subset of the rows of $Y$ dominates $Z^{\prime}$, the conjugate of $Z$.

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Examples: wide


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## Examples: not wide



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Theorem:
A wide Young diagram with at most two distinct row lengths is Latin.

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A self-conjugate wide Young diagram with at most three distinct row lengths is Latin.

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## Theorem:

If the wide partition conjecture holds for selfconjugate wide diagrams, then it holds for all wide diagrams.

## The 3-hypergraph $H(Y)$

To a Young diagram $Y$ we assign a tripartite 3-hypergraph $H(Y)$ as follows:

Sides:

$$
\begin{array}{ll}
R=R(Y)=\left\{r_{1}, r_{2}, \ldots, r_{m}\right\} & \text { (the rows) } \\
C=C(Y)=\left\{c_{1}, c_{2}, \ldots, c_{a_{1}}\right\} & \text { (the columns) } \\
S=S(Y)=\left\{s_{1}, s_{2}, \ldots, s_{a_{1}}\right\} & \text { (the symbols) }
\end{array}
$$

Edges:

$$
E(H(Y))=\left\{r_{i} c_{j} s_{k} \mid 1 \leq i \leq m, c_{j} \in C, s_{k} \in S, 1 \leq j, k \leq a_{i}\right\}
$$

## The 3-hypergraph $H(Y)$

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An edge $r_{i} c_{j} s_{k}$ in $H(Y)$ corresponds to the cell $(i, j)$ with the symbol $s_{k}$ in it.

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An edge $r_{i} c_{j} s_{k}$ in $H(Y)$ corresponds to the cell $(i, j)$ with the symbol $s_{k}$ in it.

A filling of $Y$ with symbols from $S$ corresponds to a set of $|Y|$ edges in $H(Y)$.

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A filling of $Y$ with symbols from $S$ corresponds to a set of $|Y|$ edges in $H(Y)$.

A Latin filling of $Y$ with symbols from $S$ corresponds to a set of $|Y|$ edges in $H(Y)$ no two of which share more than one vertex.

## The $k$ th matching number

Definition: (Aharoni and Zerbib*, 2020)
a $k$-matching in a hypergraph $H$ is a subset of $\mathrm{E}(H)$ in which every two edges share fewer than $k$ vertices.

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The $k$ th matching number $v^{(k)}(H)$ is the maximal size of a $k$ matching in $H$.

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WPC is equivalent to:
Wide Partition Conjecture (hypergraph version):
If $Y$ is wide, then $v^{(2)}(H(Y))=|Y|$

[^5]
## The $k$ th covering number

Definition: (Aharoni and Zerbib, 2020)
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The $k$ th covering number $\tau^{(k)}(H)$ is the minimal size of a $k$-cover in $H$.

Easy to see:

$$
\tau^{(k)}(H) \geq v^{(k)}(H)
$$

(Given a $k$-matching $M$ of maximal size $v^{(k)}(H)$, we need at least $|M| k$-sets of edges to cover its members.)

## A weak version of WPC

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If a Young diagram $Y$ is wide, then $\tau^{(2)}(H(Y))=|Y|$

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If $\tau^{(2)}(H(Y))=|Y|$, then the Young diagram $Y$ is wide
This strengthens Observation 1:
$Y$ is Latin $\Rightarrow v^{(2)}(H(Y))=|Y| \Rightarrow \tau^{(2)}(H(Y))=|Y| \Delta Y$ is wide

## A fractional version

A fractional 2-matching in a hypergraph $H$ is a function $f: E(H) \rightarrow \mathbb{R}_{\geq 0}$ subject to the constraint

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v^{(2) *}(H)=\max _{\substack{f \text { fractional } \\
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\end{gathered}
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A fractional 2-cover of a hypergraph $H$ is a function $g:\binom{V(H)}{2} \rightarrow \mathbb{R}_{\geq 0}$ subject to the constraint

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By the definition and LP duality:
$\tau^{(k)}(H) \geq \tau^{(k) *}(H)=\nu^{(k) *}(H) \geq \nu^{(k)}(H)$

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## A fractional version

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Conjecture:
If $Y$ is a wide Young diagram, then $Y$ is fractionally Latin
$v^{(2)}$ Vs. $\tau^{(2)}$

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Observation:
For a Young diagram $Y \tau(H(Y))=v(H(Y))$

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## Question:

Is it true that $\tau^{(2)}(H(Y))=v^{(2)}(H(Y))$ for every Young diagram $Y$ ?

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If the Wide Partition Conjecture is true, then the answer is YES

## M-tableaux

$a_{1}$

$m$ rows: $\quad r_{1}, r_{2}, \ldots, r_{m}$
row lengths: $a_{1}, a_{2}, \ldots, a_{m}$ respectively
$a_{1}$ columns: $c_{1}, c_{2}, \ldots, c_{a_{1}}$
M a matroid
An $M$-tableau is a Young diagram with an element of $M$ in each of its cells

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$M$ a matroid
An $M$-tableau is a Young diagram with an element of $M$ in each of its cells

## Question (Chow, Fan, Goemans, Vondrak 2003*):

Given an $M$-tableau such that the elements in each row are independent, is it possible to rearrange the elements in each row so that the elements in each column are independent?

[^6]
## Generalized Rota's Basis Conjecture

Rota's Basis Conjecture (Huang and Rota 1994):
If all the rows are of size $\operatorname{rank}(M)$ and $m \leq \operatorname{rank}(M)$, then the answer is Yes

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Generalized WPC (Chow, Fan, Goemans, Vondrak 2003):
The answer is Yes if and only if $Y$ is wide



[^0]:    * T. Y. Chow, C. K. Fan, M. X. Goemans, and J. Vondrak. Wide partitions, Latin tableaux, and Rota's basis coniecture. Adv. in Appl. Math. 31 (2003), 334-358.

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[^2]:    *R. Aharoni and S. Zerbib. A generalization of Tuza's conjecture. J. Graph Theory 94 (2020), 445-462.

[^3]:    *R. Aharoni and S. Zerbib. A generalization of Tuza's conjecture. J. Graph Theory 94 (2020), 445-462.

[^4]:    *R. Aharoni and S. Zerbib. A generalization of Tuza's conjecture. J. Graph Theory 94 (2020), 445-462.

[^5]:    *R. Aharoni and S. Zerbib. A generalization of Tuza's conjecture. J. Graph Theory 94 (2020), 445-462.

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